Security functions in mobile communication systems

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- Security demands
- Security functions of GSM
- Known attacks on GSM
- **Security functions of UMTS**
- Concepts for hiding locations of mobile users

Security deficits of existing mobile networks

• Example of security demands: Cooke, Brewster (1992)

- protection of user data
- protection of signaling information, incl. location
- user authentication, equipment verification
- fraud prevention (correct billing)

• Security deficits of GSM (selection)

- Only symmetric cryptography (algorithms no officially published)
- Weak protection of locations (against outsiders)
- No protection against insider attacks (location, message content)
- No end-to-end services (authentication, encryption)

Summary

- GSM provides protection against <u>external attacks</u> only.
- "...the designers of GSM did not aim at a level of security much higer than that of the fixed trunk network." Mouly, Pautet (1992)

Security functions of GSM

- Overview
 - Subscriber Identity Module (SIM, smart card)
 - Admission control and crypto algorithms
 - -Authentication (Mobile station \rightarrow network)
 - Challenge-Response-Authentication (A3)
 - Pseudonymization of users on the air interface
 - Temporary Mobile Subscriber Identity (TMSI)
 - Link encryption on the air interface
 - Generation of session key: A8
 - Encryption: A5



Challenge-Response-Authentication

• When initialized by the mobile network?

- Location Registration
- Location Update when changing the VLR
- Call Setup (both directions)
- Short Message Service



Challenge-Response-Authentication

Algorithm A3

- Implemented on SIM card and in Authentication Center (AuC)
- Cryptographic one way function A3:

SRES' = A3(Ki, RAND) (Ki: individual user key)

- Interfaces are standardized, cryptographic algorithm not standardized
- Specific algorithm can be selected by the network operator
 - Authentication data (RAND, SRES) are requested from AuC by the visited MSC
 - visited MSC: only compares SRES == SRES'
 - visited MSC has to trust
 home network operator



Attacks – Telephone at the expense of others

SIM cloning

 Weakness of authentication algorithm

- Interception of authentication data
 - -Eavesdropping of internal communication links
- IMSI catcher
 - -Man-in-the-middle attack on the air interface

SIM cloning

• Scope

- Telephone at the expense of others
- Described by Marc Briceno (Smart Card Developers Association), Ian
 Goldberg and Dave Wagner (both University of California in Berkeley)
- http://www.isaac.cs.berkeley.edu/isaac/gsm.html
- Attack uses a weakness of algorithm COMP128, which implements A3/A8
- SIM card (incl. PIN) must be under control of the attacker for at least 8-12 hours

• Effort

- Approx. 150.000 calculations to determine Ki (max. 128 bit)
- 6,25 calculations per second only, due to slow serial interface of SIM card

Interception of authentication data

• Scope

- Telephone at the expense of others

- Described by Ross Anderson (University of Cambridge)
- Eavesdropping of unencrypted internal transmission of authentication data (RAND, SRES) from AuC to visited MSC

Weakness

- GSM standard only describes interfaces between network components.
- They forgot the demand for internal encryption.
- Microwave links are widely used for internal linkage of network components.

No encryption of internal links



Interception of authentication data



IMSI-Catcher

• Scope

- Identities of users of a certain radio cell
- Eavesdropping of communications
- (Telephone at the expense of others)
- Man-in-the-middle attack (Masquerade)
- Weakness
 - No protection against malicious or faked network components



Universal mobile telecommunication system (UMTS)

- Security functions of UMTS ...
 - ... have been »inspired« by GSM security functions

• From GSM

- Subscriber identity confidentiality (TMSI)
- Subscriber authentication
- Radio interface encryption
- SIM card (now called USIM)
- Authentication of subscriber towards SIM by means of a PIN
- Delegation of authentication to visited network
- No need to adopt standardized authentication algorithms

Additional UMTS security features

- Enhanced UMTS authentication and key agreement mechanism
- Integrity protection of signaling information (prevents false-base-station attacks)
- New ciphering / key agreement / integrity protection algorithms
- ... and a few minor features

UMTS Security Architecture

USIM MS

RNC

VLR

SGSN

HLR

AuC

Authentication Centre





UMTS Subscriber Identity Moduleauthentication key K,Mobile Stationauthentication function f1, f2Mobile Stationkey generation function f3, f4, f5Radio Network Controllersequence number management SQNVisitor Location Reg.sequence number management SQNSG Serving NetworkHome Location Register

AV := RAND || XRES || CK || IK || AUTN

AUTN := SQN AK || AMF || MAC



Generation of authentication vectors

Authentication function in the USIM



Verify MAC == XMAC

Verify that SQN is in the correct range

Security mode setup procedure



Cipher algorithm f8

- Combination of Output Feedback mode (OFB) and counter mode
- First encryption under CK' prevents chosen plaintext attacks (initialization vector is encrypted, KM: key modifier)



Integrity algorithm f9

- ISO/IEC 9797-1 (MAC algorithm 2)
- Sender and receiver use f9
- Receiver verifies MAC == XMAC



Protection of locations

- Mobile user
 - whishes to be reachable at his current location.
 - He won't be localizable by outsiders and the network operator unless the explicitly gives his permission
- There is no mobile network that fulfills this demand.



Protection of locations

- GSM (Global System for Mobile Communication)
 - Distributed storage at location registers
 - Home Location Register (HLR)
 - Visitor Location Register (VLR)
 - Network operator has global view on location information
- Tracking of mobile users is possible



Systematic: Protection of locations

A. Trust into the mobile station only

- A.1 Broadcast method
- A.2 Group pseudonyms

B. Additional trust into a private fixed station

- B.1 Trusted address translation and broadcast
- B.2 Reduction of broadcast areas
- B.3 Explicit trustworthy storage of locations
- B.4 Temporary pseudonyms (TP method)

C. Additional trust into a trusted third party

- C.1 Trust Center
- C.2 Co-operating chips
- C.3 Mobile Communication-MIXing



• Overview: Broadcast

• No storage of locations and global paging of mobile users



• Overview: Broadcast

• No storage of locations and global paging of mobile users



• Immense costs for bandwith ...

Broadcast in general



• Overview : Trustworthy storage

Replace databases by trusted devices in the fixed network



Overview : Trustworthy storage

Replace databases by trusted devices in the fixed network



- Every location updating needs communication with trusted station.
- Question: How can we reduce cost of location updating?

Overview : Trustworthy storage

Tempory Pseudonyms (TP method)



• Can we do this without a trusted fixed station?

Overview : Mobile Communication-MIXing

Covered storage of location information



A MIX hides the communication relation between

- HLR and VLR
- VLR and location area

Implicit Addresses

- First contact: Covered Implicit Address CIA
 - Recipient publishes public encryption key c
 - Sender creates CIA := c(R,S,M)
 - Redundancy R
 - Seed S of a pseudo-random generator PRG
 - Message **M** (optional, may contain symmetric key **k**)
 - Recipient decrypts all received messages with private key d
 - Finds correct R for own messages only
- Following addressing: Open Implicit Address OIA
 - -OIA_{i+1} := PRG(i,seed) (i = 0,1,2,...)
 - Sender :
 - calculates next OIA
 - encrypts message (optional) **M** under k
 - Sends OIA, M
 - Receiver: Associative memory of all valid OIAs to recognize own messages

Broadcast method

Performance



Performance: Message lengths on the air interface

Mobile Terminated Calls

- GSM reference
- B.3 explicit trustworthy storage
- B.4 TP method
- C.2 cooperating chips



Performance: Message lengths on the air interface

Location Update

- GSM reference
- B.3 explicit trustworthy storage
- B.4 TP method
- C.2 cooperating chips



Security of mobile communication

Conclusion

- Protection of locations can be technically realized

- However, there is a demand for legal enforcement

More information

- http://www.inf.tu-dresden.de/~hf2/mobil/

